

# Generation of Signals under Temporal Constraints for CPS Testing

Benoît Barbot<sup>1</sup>, Nicolas Basset<sup>2</sup> and Thao Dang<sup>2</sup>

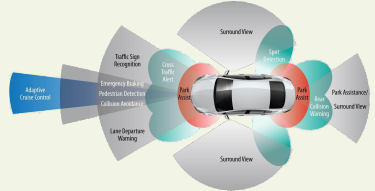
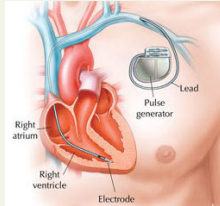
<sup>1</sup>LACL, Université Paris Est Créteil, France

<sup>2</sup>VERIMAG/CNRS, Université Grenoble Alpes, France

Vendredi 24 Mai 2019, Mefosyloma

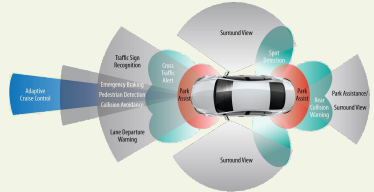
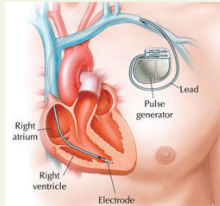
# Cyber Physical System (CPS)

## Examples



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## Examples



## System with three different interacting aspects

- ▷ A continuous deterministic part following differential equations
- ▷ A discrete part with deterministically guarded transitions
- ▷ Interactions with the environment

# Cyber Physical System Analysis

## State-space

- ▷ Continuous part difficult to analyse
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Testing is the only tractable validation technique

Require the generation of real-valued signal as input of the CPS.

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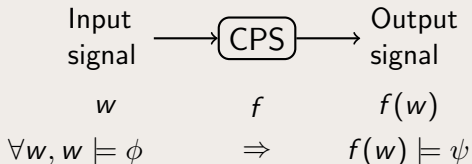
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## CPS testing problem



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### Using time automaton to model signal [Alur & Dill, 1994]

- Well suited for modelling time constraints between events.
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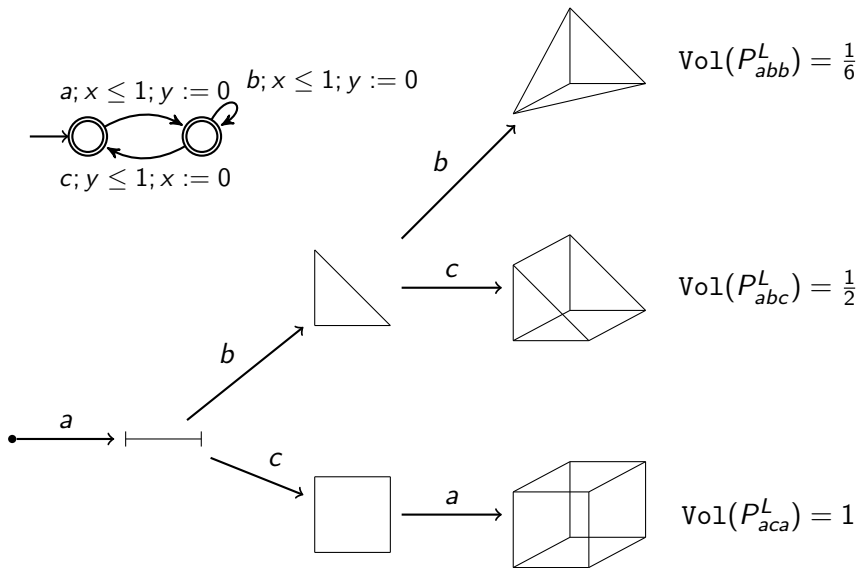
### Using time automaton to model signal [Alur & Dill, 1994]

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⇒ require to sample time words from time automaton language.

## Geometrical shape of time language

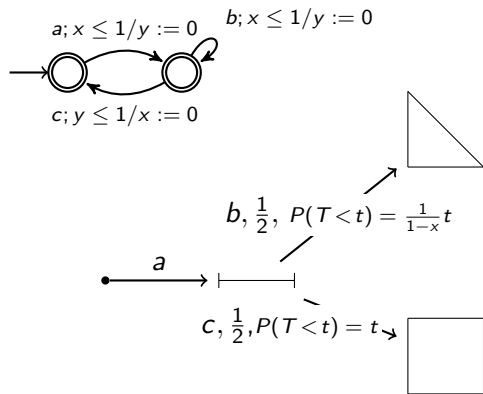
Constraints of timings along a path  $w = \text{polytope } P_w^L$



## Isotropic sampling (a "by-default" sampling)

At each step...

- ▷ assign same weight to every edge  $e$  and pick one randomly;
- ▷ assign same weight to every time  $t$  such that  $(t, e)$  can be taken.

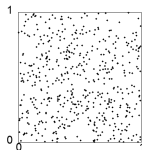
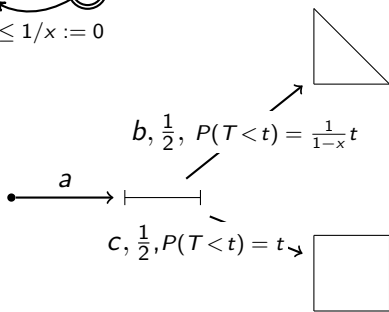
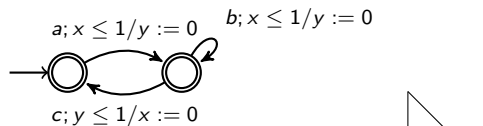




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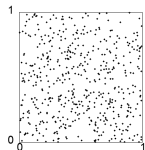
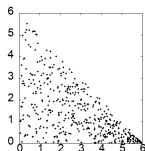
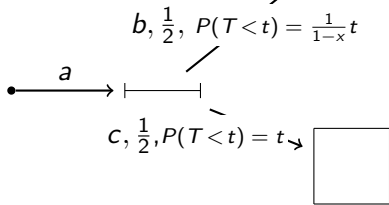
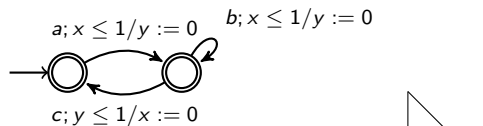
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## Polytopes and CDF

### Uniform distribution for a fixed word

Given a word  $\alpha$  in the language of the automaton; and  $\mathcal{P}$  the associated polytope. The uniform distribution over timed vector  $\mathbf{t} \in \mathcal{P}$  assigns the density probability  $\omega(\mathbf{t}) = \frac{1}{\text{Vol}(\mathcal{P})}$ .

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### Cumulative Density Function (CDF)

The CDF of the uniform distribution over  $\alpha$  is  $F : \mathcal{P} \rightarrow [0; 1]$  defines as  $F(\mathbf{t}) = \int_{\mathbf{T} < \mathbf{t}} \omega(\mathbf{T}) d\mathbf{T}$ . Moreover, the CDF  $F$  can be written as:

$$F(t_1, \dots, t_n) = F_1(t_1)F_2(t_2|t_1) \dots F_n(t_n|t_1, \dots, t_{n-1})$$

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### In QUEST16

Given a path in a TA; one can effectively compute the CDF of the uniform distribution in the form  $F_i(t_i|t_1, \dots, t_{i-1}) = \frac{\pi_i(t_1, \dots, t_i)}{\gamma_i(t_1, \dots, t_{i-1})}$  with  $\pi_i$  and  $\gamma_i$  polynomials of degree  $i$ .

## QEST16 in a nutshell

### Method

- ▷ Compute the forward reachable zone-graph of the automaton.
- ▷ Compute recursively the volume of language for each state of the zone-graph.
- ▷ The PDF  $\omega(t)$  is the normalized volume.

### Recursive definition over state of the zone-graph

$$v_0(s) = 1$$
$$v_{n+1}(s) = \sum_{\delta \in \Delta(s)} \int_{lb_\delta(s)}^{ub_\delta(s)} v_n(s_{(t,\delta)}) dt.$$

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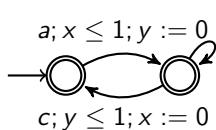
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### Shape of zone matter !

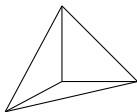
Zone-graph require additionnal splitting to ensure that the integral is simple.

# Example



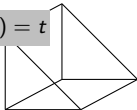
$$b, \frac{1-x}{2-x}, P(T < t) = \frac{t}{1-x}$$

$$\text{Vol}(P_{abb}^L) = \frac{1}{6}$$

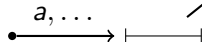


$$c, \frac{1}{2-x}, P(T < t) = t$$

$$\text{Vol}(P_{abc}^L) = \frac{1}{2}$$



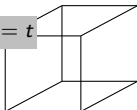
$$b, \frac{2.5-4x+1.5x^2}{3.5-4x+1.5x^2}, P(T < t) = \frac{4t-2xt+t^2}{5-8x+3x^2}$$



$$c, \frac{1}{3.5-4x+1.5x^2}, P(T < t) = t$$

$$a, \frac{1}{1}, P(T < t) = t$$

$$\text{Vol}(P_{aca}^L) = 1$$





## Sampling polytopes

Unit cube

In practice sampling methods produce floating point number in  $[0; 1]$ .

How to sample the uniform distribution ?

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## Inverse Sampling

Recalls  $F_i(t_i | t_1, \dots, t_{i-1})$  is a polynomial in  $t_i \rightarrow [0; 1]$ .

- Sample a real  $u$  in  $[0; 1]$
- Compute  $t_i$  as the root of  $F_i(t_i | t_1, \dots, t_{i-1}) - u$

*Note that  $F_i$  is a strictly increasing polynomial  $\Rightarrow$  Newton method applies.*

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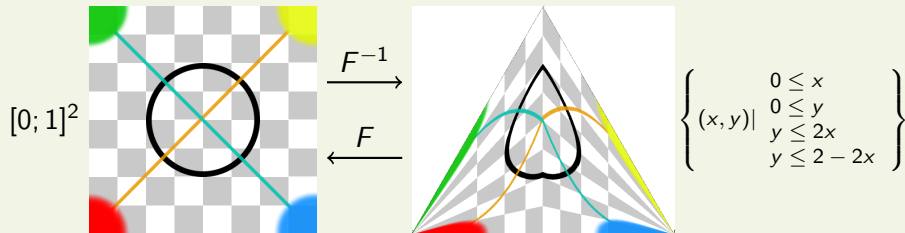
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$$F^{-1} : [0; 1]^n \rightarrow \mathcal{P}$$

- ▷ Given  $\mathbf{u}$  in  $[0; 1]^n$
- ▷ Apply inverse sampling iteratively to obtain  $\mathbf{t} \in \mathcal{P}$

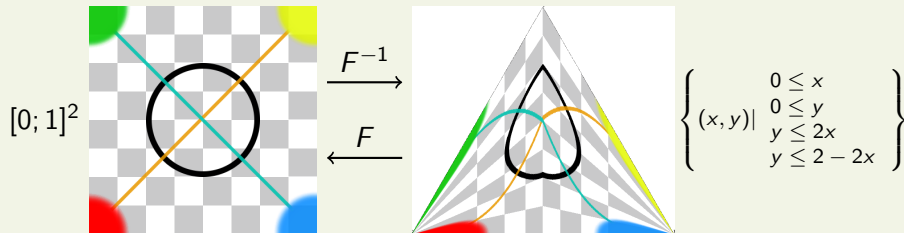
## Sampling polytopes II

Example in dimension 2



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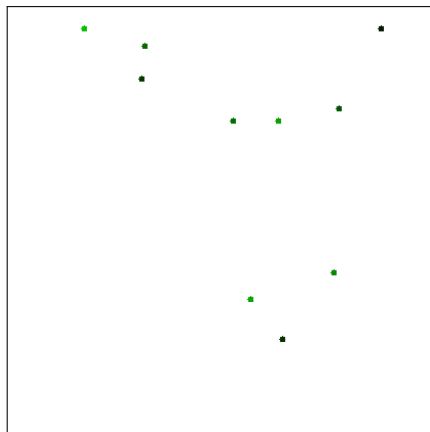


### Sampling methods

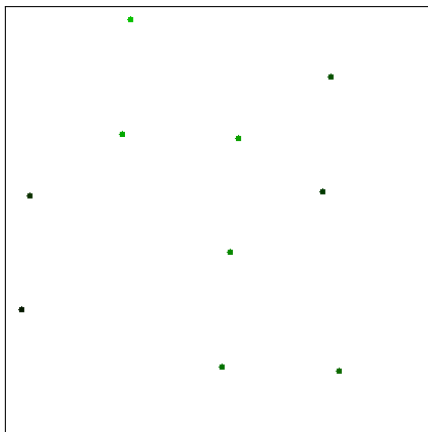
- ▷ Uniform (Pseudo) Random number
- ▷ Low discrepancy sequence

# Random vs low discrepancy sequence (Empirical)

Uniform Random



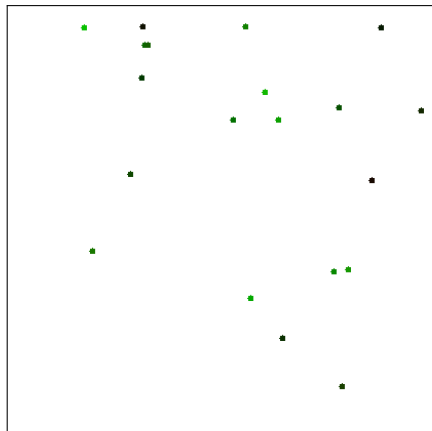
Kronecker



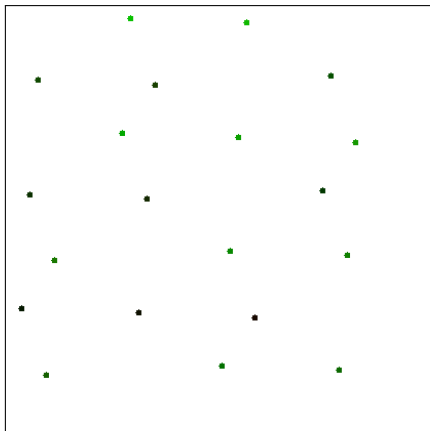
$n = 10$

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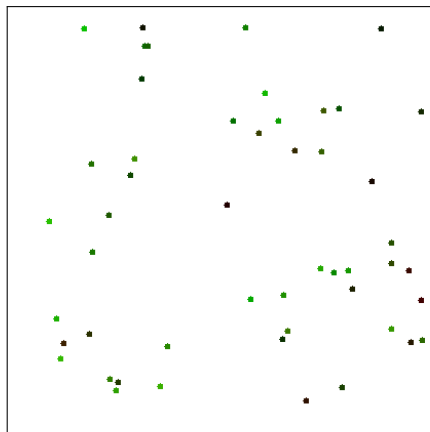
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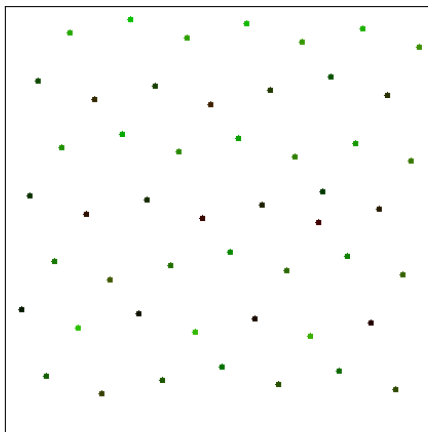
$n = 20$

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Uniform Random



Kronecker

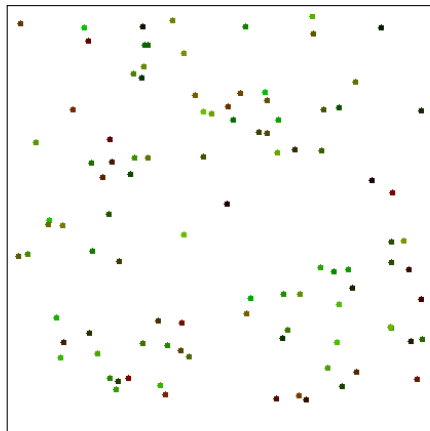


$n = 50$

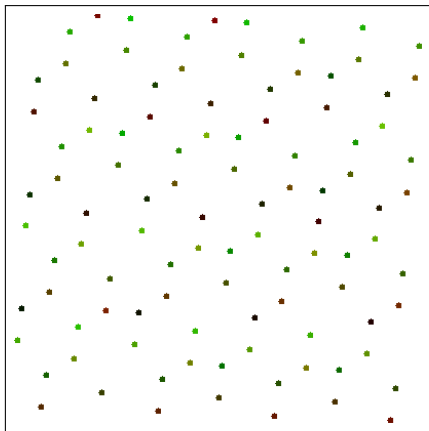


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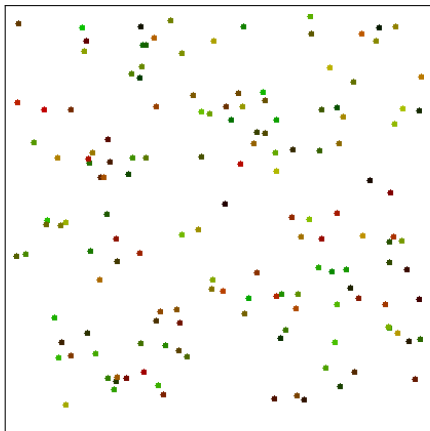
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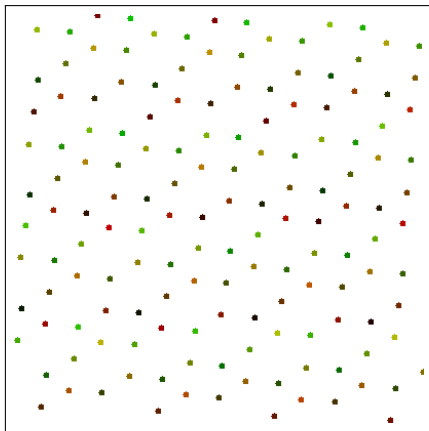
$n = 100$

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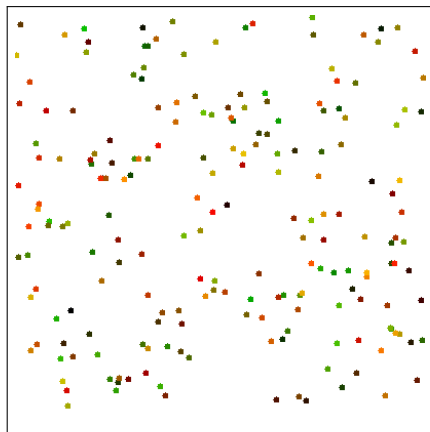
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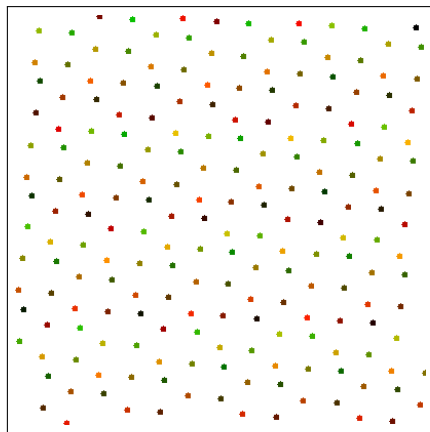
$n = 150$

# Random vs low discrepancy sequence (Empirical)

Uniform Random



Kronecker



$n = 200$

# Discrepancy

## Star Discrepancy Definition

For  $\mathbf{b} = (b_1, \dots, b_n) \in [0, 1]^n$ , we define the box  $[\mathbf{0}, \mathbf{b}] = [0, b_1] \times \dots \times [0, b_n]$ . The star discrepancy of a finite set  $S$  is defined as:

$$D_{\star}(S) = \sup_{\mathbf{b} \in [0, 1]^n} \left| \text{Vol}([\mathbf{0}, \mathbf{b}]) - \frac{|S \cap [\mathbf{0}, \mathbf{b}]|}{|S|} \right|.$$

## Random vs low discrepancy sequence (theory)

For  $g : [0, 1]^n \rightarrow [a, b]$ :

### (Pseudo) Random sequence

Guarantee given as probabilistic framing i.e. confidence interval.

ex Chernoff-Hoeffding bounds: Let  $z > 0$ , let  $1 - 2e^{-2z^2}$  be the confidence level:

$$\mathbb{P} \left( \left| \frac{1}{N} \sum_{n=1}^N g(\mathbf{p}^{(n)}) - \int_{[0,1]^n} g(\mathbf{r}) d\mathbf{r} \right| < 2z \frac{b-a}{\sqrt{N}} \right) \geq 1 - 2e^{-2z^2}$$

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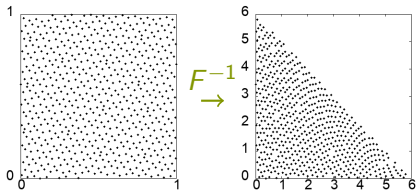
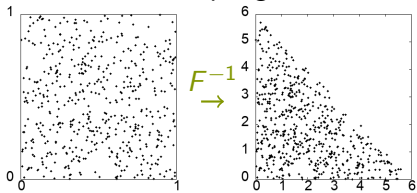
### Low Discrepancy Sequence

Guarantee given as deterministic framing i.e. interval using Koksma-Hlawka inequality.

$$\left| \frac{1}{N} \sum_{n=1}^N g(\mathbf{p}^{(n)}) - \int_{[0,1]^n} g(\mathbf{r}) d\mathbf{r} \right| \leq V^*(g)(D_*(S))^{1/n}$$

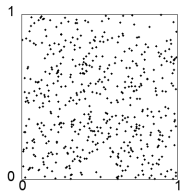
# Application to the Evaluation of Uniformity Degree

Sampling

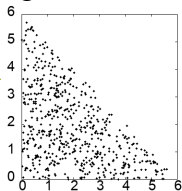


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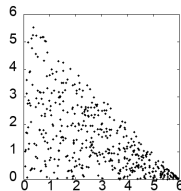
Sampling



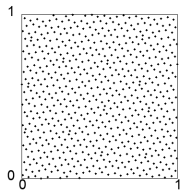
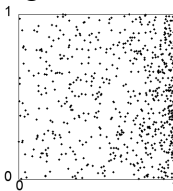
$F^{-1}$   
→



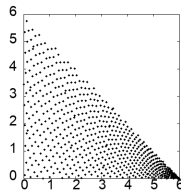
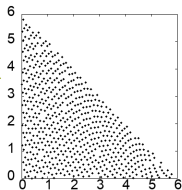
Evaluating



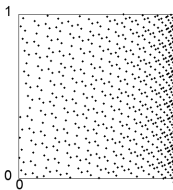
$F$   
→



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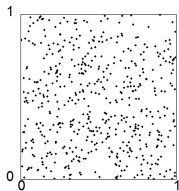
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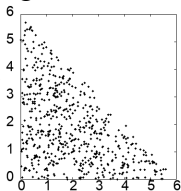


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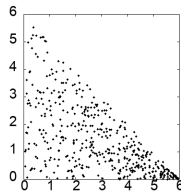
Sampling



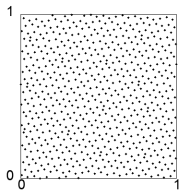
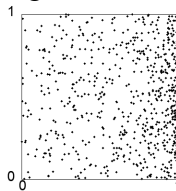
$F^{-1}$   
→



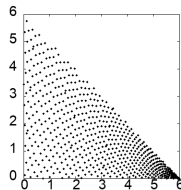
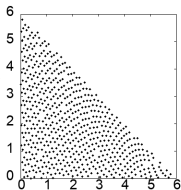
Evaluating



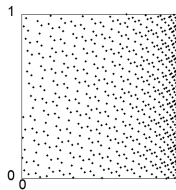
$F$   
→



$F^{-1}$   
→



$F$   
→

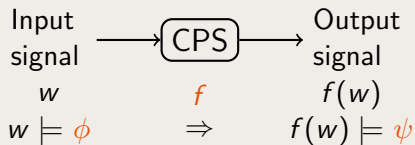


## Kolmogorov-Smirnov test

- ▷ Quantify the distance between two distributions.
- ▷ When apply between an empirical distribution  $S$  and the uniform one, equivalent to  $D_*(S)$ .

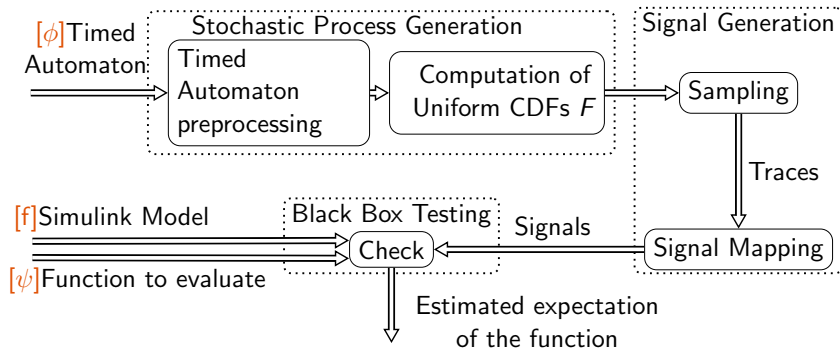
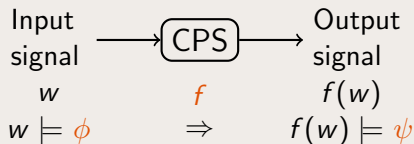
# Application To CPS Testing

## CPS Testing Abstract



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## CPS Testing Abstract



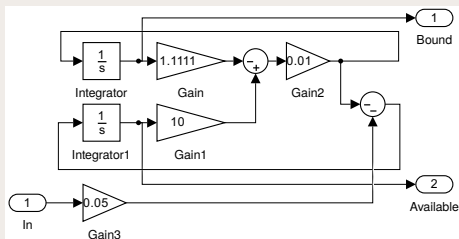
## Example: KiBaM System

### A CPS with a controller and a battery

$a : x < \tau_3 \wedge y > \tau_2; x := 0$

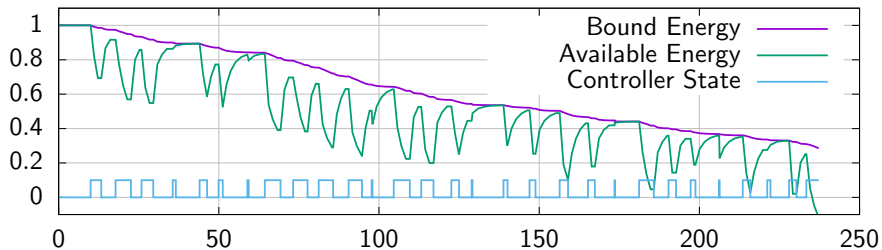


$b : x < \tau_1; y := 0$

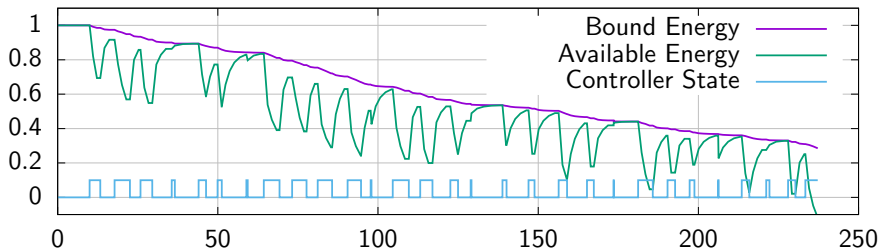


- The controller is switched on at least every  $\tau_3$  time unit.
- The controller consume energy when switch on.
- The battery self-recharge on low load.

## Example: KiBaM System Results



## Example: KiBaM System Results



$(Available > 0) U^{\leq 150} T$

CDF computation: 10s

Number of trajectories: 1,000,000

Uniform Random: Falsifies 53 trajectories, 1400s of simulations

Low discrepancy sequence: Falsifies 56 trajectories, 1600s of simulations

# Stability of a $\Sigma\Delta$ Modulator

## Characteristics

- Analog to digital converter;
- Mixing discrete-time and continuous-time components;
- Subject to saturation;

⇒ treated as a black box.

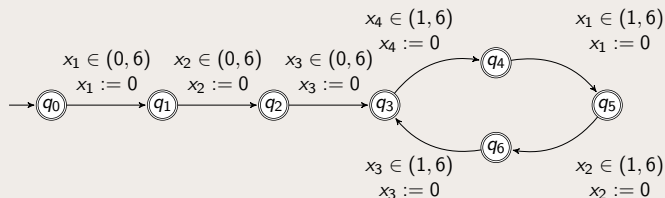
# Stability of a $\Sigma\Delta$ Modulator

## Characteristics

- Analog to digital converter;
- Mixing discrete-time and continuous-time components;
- Subject to saturation;

⇒ treated as a black box.

## Input Signals



- A pseudo-periodic signal.
- Signals are linear interpolation based on location.
- CDF computation: 30s



## Stability of a $\Sigma\Delta$ Modulator II

$(\neg \text{saturation}) \cup_{\leq \text{simtime}} \top$

- Several batches with a scaling parameter  $\kappa$  for frequencies
- 100 trajectories per batch, simulated in 1 minute
- Several test with scaling parameter for frequencies
  - $\kappa \geq 0.8 \times 10^{-7} \Rightarrow$  saturation detected with both methods
  - $\kappa = 0.6 \times 10^{-7} \Rightarrow$  saturation detected with low discrepancy sequence
  - $\kappa \leq 0.5 \times 10^{-7} \Rightarrow$  no saturation detected

# Conclusion and Perspective

## Conclusion

- ▷ Combining uniform word generation and low-discrepancy sampling;
- ▷ Validation of complex CPS system.

# Conclusion and Perspective

## Conclusion

- ▷ Combining uniform word generation and low-discrepancy sampling;
- ▷ Validation of complex CPS system.

## Perspective

- Replacing automaton specification by a specification logic;
- Better sampling of the signal value space;
- An easy to use, self-contained implementation;
- Computing star discrepancy efficiently.